Lecture 1: an invitation to Priestley duality

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TACL 2015 School Campus of Salerno (Fisciano) 15–19 June 2015

Outline

Bounded distributive lattices

Priestley duality for finite distributive lattices

Priestley duality via homsets

Priestley duality for infinite distributive lattices

Examples of Priestley spaces

The translation industry: restricted Priestley duals

Useful facts about Priestley spaces

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Concrete examples of bounded distributive lattices

1. The two-element chain

$$\underline{\boldsymbol{2}} = \langle \{0,1\}; \vee, \wedge, 0,1 \rangle.$$



2. All subsets of a set S:

$$\langle \mathscr{P}(S); \cup, \cap, \varnothing, S \rangle.$$

3. Finite or cofinite subsets of \mathbb{N} :

$$\langle \wp_{FC}(\mathbb{N}); \cup, \cap, \varnothing, \mathbb{N} \rangle.$$

4. Open subsets of a topological space X:

$$\langle \mathcal{O}(\mathbf{X}); \cup, \cap, \varnothing, X \rangle.$$

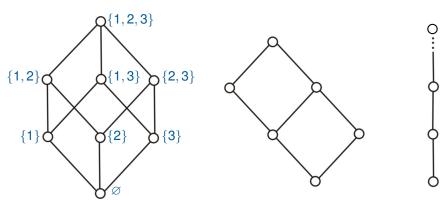
More examples of bounded distributive lattices

- 5. $\langle \{T, F\}; \text{ or, and, } F, T \rangle$.
- 6. ⟨N ∪ {0}; lcm, gcd, 1, 0⟩.
 (Use the fact that, lcm(m, n) · gcd(m, n) = mn, for all m, n ∈ N ∪ {0}, and that a lattice is distributive iff it satisfies x ∨ z = y ∨ z & x ∧ z = y ∧ z ⇒ x = y.)
- 7. Subgroups of a cyclic group **G**,

$$\langle \mathsf{Sub}(\mathbf{G}); \, \vee, \cap, \{e\}, G \rangle$$
, where $H \vee K := \mathsf{sg}_{\mathbf{G}}(H \cup K)$.

Drawing distributive lattices

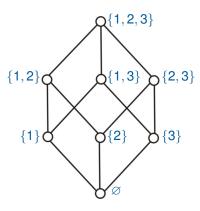
Any distributive lattice $\langle L; \vee, \wedge, 0, 1 \rangle$ has a natural order corresponding to set inclusion: $a \leq b \iff a \vee b = b$.

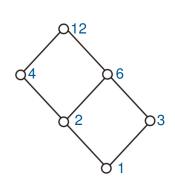


- ∨ union
- ∧ intersection
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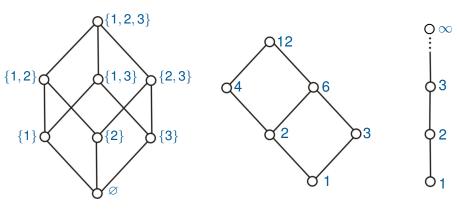


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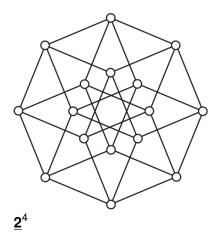
- ∨ lcm
- ∧ gcd

/ max

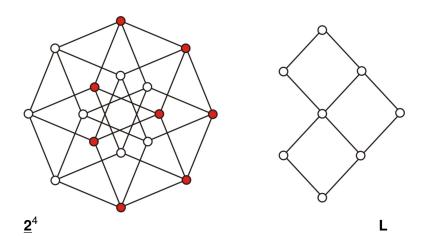
∧ min

< usual

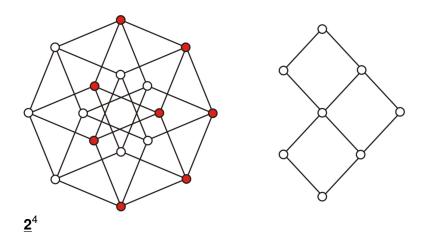
A sublattice ${\bf L}$ of ${\bf \underline{2}}^4$



A sublattice ${\bf L}$ of ${\bf \underline{2}}^4$



A sublattice L of 2⁴



Note: Every distributive lattice embeds into $\underline{\mathbf{2}}^{S}$, for some set S.

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Representing finite distributive lattices

Birkhoff's representation for a finite distributive lattice **L**

Let **L** be a finite distributive lattice.

L is isomorphic to the collection $\mathcal{O}(\mathbf{P})$ of all down-sets of an ordered set $\mathbf{P} = \langle P; \leqslant \rangle$, under union, intersection, \varnothing and P.

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Theorem [G. Birkhoff]

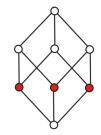
Let **L** be a finite distributive lattice and let **P** be a finite ordered set. Then

- ▶ **L** is isomorphic to $\mathcal{O}(\mathcal{J}(\mathbf{L}))$, and
- ▶ **P** is isomorphic to $\mathcal{J}(\mathcal{O}(\mathbf{P}))$.

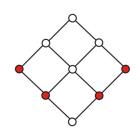
More examples

Distributive lattice

 $\textbf{L}\cong\mathcal{O}(\mathcal{J}(\textbf{L}))$







Ordered set

$$\mathsf{P}\cong\mathcal{J}(\mathcal{O}(\mathsf{P}))$$



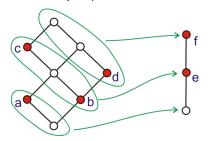


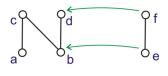


Duality for finite distributive lattices

The classes of

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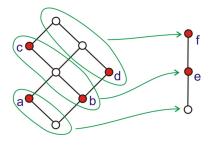


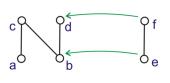


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 $\begin{array}{cccc} \text{surjections} & \longleftrightarrow & \text{embeddings} \\ \text{embeddings} & \longleftrightarrow & \text{surjections} \\ \text{products} & \longleftrightarrow & \text{disjoint unions} \end{array}$

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Duals of finite bounded distributive lattices

Let $\mathbf{L} = \langle L; \vee, \wedge, 0, 1 \rangle$ be a finite bounded distributive lattice. We can define its dual $D(\mathbf{L})$ to be either

- J(L) the ordered set of join-irreducible elements of L
 or
 - ▶ $\mathcal{D}(\mathbf{L}, \mathbf{\underline{2}})$ the ordered set of $\{0, 1\}$ -homomorphisms from \mathbf{L} to the two-element bounded lattice $\mathbf{\underline{2}} = \langle \{0, 1\}; \lor, \land, 0, 1 \rangle$.

Here $\mathfrak D$ denotes the category of bounded distributive lattices.

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In fact, we have the following dual order-isomorphism:

$$\mathcal{J}(\mathsf{L}) \cong^{\partial} \mathfrak{D}(\mathsf{L}, \underline{\mathbf{2}}).$$

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Duals of finite ordered sets

Let $\mathbf{P} = \langle P; \leqslant \rangle$ be a finite ordered set.

We can define its dual $E(\mathbf{P})$ to be either

 $ightharpoonup \mathcal{O}(\mathbf{P})$ – the lattice of down-sets (= order ideals) of \mathbf{P}

or

▶ $\mathcal{P}(\mathbf{P}, \mathbf{Z})$ – the lattice of order-preserving maps from \mathbf{P} to the two-element ordered set $\mathbf{Z} = \langle \{0, 1\}; \leqslant \rangle$.

Here \mathcal{P} denotes the category of ordered sets.

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Duals of morphisms

Let **L** and **K** be a finite distributive lattices and let **P** and **Q** be a finite ordered sets.

There is a bijection between the {0,1}-homomorphisms from L to K and the order-preserving maps from D(K) to D(L). Given f: L → K, we define

$$\varphi \colon \mathcal{J}(\mathbf{K}) \to \mathcal{J}(\mathbf{L}) \text{ by } \varphi(x) := \min(f^{-1}(\uparrow x)),$$

 $\varphi \colon \mathcal{D}(\mathbf{K}, \underline{\mathbf{2}}) \to \mathcal{D}(\mathbf{L}, \underline{\mathbf{2}}) \text{ by } \varphi(x) := x \circ f.$

We denote the map φ by D(f).

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► There is a bijection between the order-preserving maps from $\bf P$ to $\bf Q$ and the $\{0,1\}$ -homomorphisms from $E(\bf Q)$ to $E(\bf P)$. Given $\varphi \colon \bf P \to \bf Q$, we define

$$f \colon \mathcal{O}(\mathbf{Q}) \to \mathcal{O}(\mathbf{P}) \text{ by } f(A) := \varphi^{-1}(A),$$

 $f \colon \mathcal{P}(\mathbf{Q}, \mathbf{Z}) \to \mathcal{P}(\mathbf{P}, \mathbf{Z}) \text{ by } f(\alpha) := \alpha \circ \varphi.$

We denote the map f by $E(\varphi)$.

The duality at the finite level

- ▶ $\mathbf{2} = \langle \{0, 1\}; \vee, \wedge, 0, 1 \rangle$ is the two-element lattice,
- ▶ $\mathbf{2} = \langle \{0, 1\}; \leqslant \rangle$ is the two-element ordered set with $0 \leqslant 1$.

Define either

$$D(\mathbf{L}) := \mathcal{J}(\mathbf{L})$$
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$$D(\mathbf{L}) := \mathfrak{D}(\mathbf{L}, \underline{\mathbf{2}}) \leqslant \underline{\mathbf{2}}^{L}$$
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Theorem [G. Birkhoff, H. A. Priestley]

Every finite distributive lattice is encoded by an ordered set:

$$L \cong ED(L)$$
 and $P \cong DE(P)$,

for each finite distributive lattice L and finite ordered set P.

Indeed, the categories of finite bounded distributive lattices and finite ordered sets are dually equivalent.

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The finite-cofinite lattice $\mathbf{L} = \langle \wp_{FC}(\mathbb{N}); \cup, \cap, \varnothing, \mathbb{N} \rangle$ cannot be obtained as the down-sets of an ordered set.

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- ▶ But $\wp_{FC}(\mathbb{N})$ is countable.

But it can be obtained as the clopen down-sets of a topological ordered set.

O O O O
$$\cdots$$
 O 1 2 3 4 5 \propto

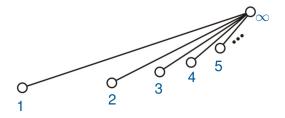
More examples

Distributive lattice:

All finite subsets of $\mathbb{N},$ as well as \mathbb{N} itself,

$$\langle \wp_{\mathsf{fin}}(\mathbb{N}) \cup \{\mathbb{N}\}; \cup, \cap, \varnothing, \mathbb{N} \rangle.$$

Topological ordered set:

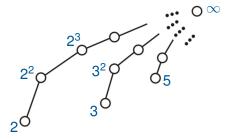


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▶ $\underline{\mathbf{2}} = \langle \{0, 1\}; \vee, \wedge, 0, 1 \rangle$ is the two-element bounded lattice.

In general, we need to endow the dual $D(\mathbf{L})$ of a bounded distributive lattice \mathbf{L} with a topology. This is easy if we define the dual of \mathbf{L} to be $D(\mathbf{L}) := \mathcal{D}(\mathbf{L}, \underline{\mathbf{2}})$.

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 $\mathfrak{D}(\mathbf{L},\underline{\mathbf{2}})$ is a topologically closed subset of $\underline{\mathbf{2}}^{L}$ (easy exercise).

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Then $D(\mathbf{L}) := \mathfrak{D}(\mathbf{L}, \underline{\mathbf{2}})$ inherits its order and topology from $\underline{\mathbf{2}}^L$.

 $\mathfrak{D}(\mathbf{L},\underline{\mathbf{2}})$ is a topologically closed subset of $\underline{\mathbf{2}}^{L}$ (easy exercise).

Hence D(L) is a compact ordered topological space.

Priestley spaces

The ordered space $D(\mathbf{L}) := \mathcal{D}(\mathbf{L}, \mathbf{\underline{2}}) \leqslant \mathbf{\underline{2}}^L$ is more than a compact ordered space. It is a Priestley space.

A topological structure $\mathbf{X} = \langle X; \leqslant, \mathfrak{T} \rangle$ is a Priestley space if

- ▶ $\langle X; \leq, \rangle$ is an ordered set,
- ightharpoonup T is a compact topology on X, and
- ▶ for all $x, y \in X$ with $x \nleq y$, there is a clopen down-set A of X such that $x \notin A$ and $y \in A$.

The category of Priestley spaces is denoted by \mathcal{P} .

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A topological structure $\mathbf{X} = \langle X; \leqslant, \mathfrak{T} \rangle$ is a Priestley space if

- ▶ $\langle X; \leq, \rangle$ is an ordered set,
- ▶ \mathfrak{T} is a compact topology on X, and
- ▶ for all $x, y \in X$ with $x \nleq y$, there is a clopen down-set A of X such that $x \notin A$ and $y \in A$.

The category of Priestley spaces is denoted by \mathcal{P} . The following result is very easy to prove.

Lemma

- ▶ $D(\mathbf{L}) = \mathcal{D}(\mathbf{L}, \mathbf{\underline{2}}) \leq \mathbf{\underline{2}}^{L}$ is a Priestley space, for every bounded distributive lattice.
- ▶ $E(X) = \mathcal{P}(X, \mathbf{2}) \leq \mathbf{2}^{X}$ is a bounded distributive lattice, for every Priestley space X.

We now have functors $D \colon \mathfrak{D} \to \mathfrak{P}$ and $E \colon \mathfrak{P} \to \mathfrak{D}$ given by

$$D(\mathbf{L}) = \mathfrak{D}(\mathbf{L}, \underline{\mathbf{2}}) \leqslant \underline{\mathbf{2}}^{L}$$
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 and $E(\mathbf{X}) = \mathfrak{P}(\mathbf{X}, \underline{\mathbf{2}}) \leqslant \underline{\mathbf{2}}^{X}$.

D and E are defined on morphisms via composition exactly as they were in the finite case:

b given $f: \mathbf{L} \to \mathbf{K}$, we define

$$D(f) \colon \mathfrak{D}(\mathbf{K}, \underline{\mathbf{2}}) \to \mathfrak{D}(\mathbf{L}, \underline{\mathbf{2}}) \text{ by } D(f)(x) := x \circ f;$$

• given $\varphi \colon \mathbf{X} \to \mathbf{Y}$, we define

$$E(\varphi) \colon \mathcal{P}(\mathbf{Y}, \mathbf{2}) \to \mathcal{P}(\mathbf{X}, \mathbf{2}) \text{ by } E(\varphi)(\alpha) := \alpha \circ \varphi.$$

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Let $\mathbf{L} \in \mathcal{D}$ and let $\mathbf{X} \in \mathcal{P}$. There are natural maps

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Priestley duality tells us that these maps are isomorphisms (in \mathcal{D} or \mathcal{P} , as appropriate).

Priestley duality

Theorem (Priestley duality)

- ▶ The functors $D: \mathcal{D} \to \mathcal{P}$ and $E: \mathcal{P} \to \mathcal{D}$ give a dual category equivalence between \mathcal{D} and \mathcal{P} .
- In particular, e_L: L → ED(L) and ε_X: X → DE(X) are isomorphisms for all L ∈ D and X ∈ P.

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Some ordered spaces

The following figure comes from Chapter 11 of Davey and Priestley: Introduction to Lattices and Order.

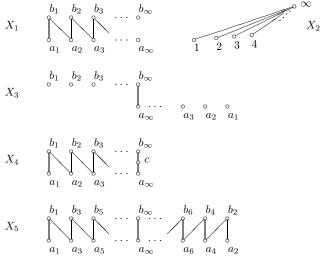


Figure 11.5

▶ If $\mathbf{X} = \langle X; \leq, \mathfrak{T} \rangle$ is a Priestley space, then \leq is a topologically closed subset of $X \times X$.

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▶ Since $\mathcal{P}(\mathbf{X}, \mathbf{Z})$ is isomorphic to the lattice $\mathcal{U}^{\mathsf{T}}(\mathbf{X})$ of clopen up-sets of \mathbf{X} , it is common to define the dual $E(\mathbf{X})$ of a Priestley space \mathbf{X} to be $\mathcal{U}^{\mathsf{T}}(\mathbf{X})$.

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- Hence when translating properties of distributive lattices into properties of Priestley spaces, it is common to use clopen up-sets (or their complements, i.e., clopen down-sets).

Examples: p-algebras

Let **X** and **Y** be a Priestley spaces and let $\varphi \colon \mathbf{X} \to \mathbf{Y}$ be continuous and order-preserving.

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- X is the dual of a distributive p-algebra, and called a p-space, iff
 - ▶ $\downarrow U$ is clopen, for every clopen up-set U;

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$$U^* = X \setminus U$$
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- If X and Y are p-spaces, then φ: X → Y is the dual of a p-algebra homomorphism iff
 - $\varphi(\max(x)) = \max(\varphi(x))$, for all $x \in X$.

(Here max(z) denotes the set of maximal elements in $\uparrow z$.)

Examples: Heyting algebras

Let **X** and **Y** be a Priestley spaces and let $\varphi : \mathbf{X} \to \mathbf{Y}$ be continuous and order-preserving.

- ► X is the dual of a Heyting algebra, and called a Heyting-space (or Esakia space), iff
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Examples: Ockham algebras

A = ⟨A; ∨, ∧, g, 0, 1⟩ is an Ockham algebra if A[♭] := ⟨A; ∨, ∧, 0, 1⟩ is a bounded distributive lattice and g satisfies De Morgan's laws and is Boolean complement on {0, 1}; in symbols,

$$g(a \lor b) = g(a) \land g(b), \ g(a \land b) = g(a) \lor g(b), \ g(0) = 1, \ g(1) = 0,$$

i.e., g is a lattice-dual endomorphism of \mathbf{A}^{\flat} .

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Thus X = ⟨X; ĝ, ≤, T⟩ will be the restricted Priestley dual of an Ockham algebra, known as an Ockham space, if X^b := ⟨X; ≤, T⟩ is a Priestley space and ĝ: X → X is an order-dual endomorphism of X^b.

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- ▶ If **X** and **Y** are Ockham spaces, then $\varphi \colon \mathbf{X} \to \mathbf{Y}$ is the dual of an Ockham algebra homomorphism if it is continuous, order-preserving and preserves the unary operation \widehat{g} .

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Prove each of the following claims. The order-theoretic dual of each statement is also true.

Let $\mathbf{X} = \langle X; \leq, \mathfrak{T} \rangle$ be a Priestley space.

- (1) The set $\downarrow Y := \{x \in X \mid (\exists y \in Y) \ x \leqslant y\}$ is closed in **X** provided *Y* is closed in **X**. In particular, $\downarrow y$ is closed in **X**, for all $y \in X$.
- (2) Every up-directed subset of **X** has a least upper bound in **X**.
- (3) The set Min(X) of minimal elements of X is non-empty.
- (4) Let Y and Z be disjoint closed subsets of X such that Y is a down-set and Z is an up-set. Then there is a clopen down-set U with Y ⊂ U and U ∩ Z = Ø.
- (5) Y is a closed down-set in **X** if and only if Y is an intersection of clopen down-sets.