

On One Embedding of Heyting Algebras

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Kuznetsov's Remark on an Embedding

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ON THE PROPOSITIONAL CALCULUS OF INTUITIONISTIC PROVABILITY

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A. V. KUZNETSOV

A. Yu. Muravitskiĭ drew my attention to:

COROLLARY 2. *Every pseudo-Boolean algebra \mathfrak{A} is a subalgebra of some Δ -enrichable pseudo-Boolean algebra in the variety generated by \mathfrak{A} .*

Enrichable Element, Enrichable Heyting Algebra

Definition: enrichable element, enrichable Heyting algebra

An element a of Heyting algebra \mathfrak{A} is said to be **enriched** by an element a^* (or a^* **enriches** a) if the following conditions are fulfilled:

- (1) $a \leq a^*$,
- (2) $a^* \rightarrow a \leq a$,
- (3) $a^* \leq x \vee (x \rightarrow a)$, for any $x \in |\mathfrak{A}|$.

An element $a \in |\mathfrak{A}|$ is **enrichable** if there exists an element $b \in |\mathfrak{A}|$ that enriches a . A Heyting algebra is **enrichable** if each element of it is enrichable.

Note

Given a Heyting algebra \mathfrak{A} and $a \in |\mathfrak{A}|$, it has been observed that if a is enriched by b and by c , then $b = c$

Kuznetsov's Theorem

Two propositional languages: \mathcal{L}_a and \mathcal{L}_m

- The propositional variables — $Var = \{p_0, p_1, \dots\}$
- \mathcal{L}_a -formulas $(A, B, \dots) := A \in Var | (A \wedge B) | (A \vee B) | (A \rightarrow B) | \neg A$
- \mathcal{L}_m -formulas $(\alpha, \beta, \dots) := \alpha \in Var | (\alpha \wedge \beta) | (\alpha \vee \beta) | (\alpha \rightarrow \beta) | \neg \alpha | \Box \alpha$
- **Int** denotes intuitionistic propositional logic in \mathcal{L}_a
- **KM** denotes proof-intuitionistic logic in \mathcal{L}_m

Kuznetsov's Theorem

$$\mathbf{KM} + A \vdash B \iff \mathbf{Int} + A \vdash B.$$

Kuznetsov's Corollary 1

Every variety of Heyting algebras is generated by its enrichable members.

The Existence of embedding derived from the Kuznetsov Theorem

Keznetsov's Corollary 2

Any Heyting algebra is embedded into an enrichable one so that these algebras separately generate one and the same variety.

Proof [Muravitsky 1988]

Step 1: Let us take a Heyting algebra \mathfrak{A} and let K be the enrichable Heyting algebras of $\text{HSP}(\mathfrak{A})$.

Step 2: $\mathfrak{A} \in \text{HSP}(K)$ [Kuznetsov's Corollary 1]

Step 3: $\mathfrak{A} \in \text{SHP}(K)$ [since any Heyting algebra has CEP]

Step 4: "algebra ... is enrichable" is a property that can be expressed by a conjunction of Horn positive formulas. Hence each of these formulas is stable under formation of direct products and homomorphic images. Therefore, $\text{HP}(K) \subseteq K$.

Step 5: Conclusion: \mathfrak{A} is embedded into an enrichable Heyting algebra \mathfrak{B} such that $\text{HSP}(\mathfrak{A}) = \text{HSP}(\mathfrak{B})$.

Question: How intimately are \mathfrak{A} and \mathfrak{B} related to one another? Which properties of \mathfrak{A} are preserved in \mathfrak{B} ?

The Embedding

Tools for the embedding: Given a Heyting algebra \mathfrak{A} , let

- $\mu_{\mathfrak{A}}$ be the partially ordered set of the prime filters of the algebra \mathfrak{A} , arranged by set inclusion;
- $\mathcal{H}(\mathfrak{A})$ be the Heyting algebra of all upward sets of $\mu_{\mathfrak{A}}$;
- $h : \mathfrak{A} \rightarrow \mathcal{H}(\mathfrak{A})$ be a Stone embedding.
- Next we define for any $X \in \mathcal{H}(\mathfrak{A})$,

$$\Delta X = \{F \mid F \in \mu_{\mathfrak{A}}, (\forall F' \in \mu_{\mathfrak{A}})(F \subset F' \Rightarrow F' \in X)\};$$

in particular,

$$\Delta h(x) = \{F \mid F \in \mu_{\mathfrak{A}}, (\forall F')(F \subset F' \Rightarrow x \in F')\}.$$

- Let $\mathcal{B}^{\Delta}(\mathfrak{A})$ be the subalgebra of $\mathcal{H}(\mathfrak{A})$, generated by the set $\{h(x) \mid x \in \mathfrak{A}\} \cup \{\Delta h(x) \mid x \in \mathfrak{A}\}$.

Definition of the embedding

Given a Heyting algebra \mathfrak{A} , we define the denumerable sequence of algebras as follows:

$$\begin{aligned}\mathfrak{A}_0 &= \mathfrak{A}, \\ \mathfrak{A}_{i+1} &= \mathcal{B}^\Delta(\mathfrak{A}_i) \quad (i < \omega).\end{aligned}$$

Along with the sequence $\{\mathfrak{A}_i\}_{i < \omega}$, we also have the embeddings:

$$\begin{aligned}\varphi_{ii} &: \mathfrak{A}_i \rightarrow \mathfrak{A}_i, i < \omega, \quad (\text{the identity embedding of } \mathfrak{A}_i) \\ \varphi_{i(i+1)} &: \mathfrak{A}_i \rightarrow \mathfrak{A}_{i+1}, i < \omega, \quad (\text{Stone embedding } h : \mathfrak{A}_i \rightarrow \mathcal{B}^\Delta(\mathfrak{A}_i)) \\ \varphi_{ij} &= \varphi_{i(i+1)} \circ \varphi_{(i+1)(i+2)} \circ \dots \circ \varphi_{(j-1)j}, \text{ where } i < j.\end{aligned}$$

Thus the sequence $\{\mathfrak{A}_i\}_{i < \omega}$ along with the embeddings φ_{ij} , $i \leq j$, form a direct family. Let $\vec{\mathfrak{A}}$ be the direct limit of this family.

Properties of the embedding

- $\vec{\mathfrak{A}}$ is an enrichable Heyting algebra.
- Each \mathfrak{A}_i is embedded into $\vec{\mathfrak{A}}$.
- If \mathfrak{A} is finite, then \mathfrak{A} and $\vec{\mathfrak{A}}$ are isomorphic.
- If \mathfrak{A} is countable, then $\vec{\mathfrak{A}}$ is also countable.
- If \mathfrak{A} is subdirectly irreducible, so is $\vec{\mathfrak{A}}$.

The main question: Is it true that $\text{HSP}(\mathfrak{A}) = \text{HSP}(\vec{\mathfrak{A}})$?

Remarks:

- 1 It must be clear that the answer to this question is affirmative if $\text{HSP}(\mathfrak{A}) = \text{HSP}(\mathfrak{A}_a^\Delta)$, where \mathfrak{A}_a^Δ is a subalgebra of $\mathcal{H}(\mathfrak{A})$ generated by $\{h(x) \mid x \in |\mathfrak{A}|\} \cup \{\Delta h(a)\}$, where a is any fixed element of \mathfrak{A} .
- 2 Since $\Delta h(a)$ enriches $h(a)$ in \mathfrak{A}_a^Δ , we have to focus on enrichment of one element of a given Heyting algebra in general setting.

Localization of Enrichment

Definition: \mathcal{E} -pair and relation \mathcal{E}

Given an algebra \mathfrak{A} and $a, a^* \in \mathfrak{A}$, (a, a^*) is an \mathcal{E} -pair (in/of \mathfrak{A}) if a is enriched by a^* in \mathfrak{A} . Then, we define:

$$\mathcal{E}_{\mathfrak{A}} = \{(a, a^*) \mid (a, a^*) \text{ is an } \mathcal{E}\text{-pair in } \mathfrak{A}\}.$$

Definition: \sim -negation, \sim -expansion

A unary operation $\sim x$ in a Heyting algebra is called **tilde-negation** (or **\sim -negation** for short) if the following identities hold:

$$\begin{array}{ll} (a) & x \rightarrow y \leq \sim y \rightarrow \sim x; & (b) & x \wedge \sim x \leq \sim \mathbf{1}; \\ (c) & \sim \mathbf{0} \leq x \vee \sim x; & (d) & \sim \mathbf{0} \rightarrow \sim \mathbf{1} \leq \sim \mathbf{1}. \end{array}$$

A Heyting \mathfrak{A} with a \sim -negation is called a **\sim -expansion** of \mathfrak{A} and denoted by (\mathfrak{A}, \sim) . The abstract class of all \sim -expansions is denoted by K .

Consequences from the two last definitions

Proposition 1

Class K is a variety.

Proposition 2

Given a \sim -negation, $(\sim\mathbf{1}, \sim\mathbf{0})$ is an \mathcal{E} -pair.

Proposition 3

Given a Heyting algebra \mathfrak{A} , if (a, a^*) is an \mathcal{E} -pair, then the operation

$$\sim x = (x \rightarrow a) \wedge a^*$$

is a \sim -negation in \mathfrak{A} so that $a = \sim\mathbf{1}$ and $a^* = \sim\mathbf{0}$.

Definition: τ -expansion, $\tau\sim$ -expansion, τ -reduct

Let \mathfrak{A} be a Heyting algebra. We enrich the signature of \mathfrak{A} with a nullary operation τ and call $\mathfrak{A}_\tau = (\mathfrak{A}, \tau)$ a τ -**expansion** of \mathfrak{A} . The \sim -expansion of \mathfrak{A}_τ satisfying in addition to the \sim -identities the identity $\sim\mathbf{1} = \tau$ is called a $\tau\sim$ -**expansion** (of \mathfrak{A}). Accordingly, the equational class of all $\tau\sim$ -expansions is denoted by K_τ . If $(\mathfrak{A}_\tau, \sim)$ is a $\tau\sim$ -expansion, we call \mathfrak{A}_τ a τ -**reduct** of the former.

Definition: packing, $\tau\sim$ -embrace, relation \triangleleft

Suppose \mathfrak{A}_τ is a subalgebra of \mathfrak{B}_τ and $(\mathfrak{B}_\tau, \sim)$ is a $\tau\sim$ -expansion generated by $|\mathfrak{A}|$. Then we say that \mathfrak{A}_τ is **packed in** \mathfrak{B}_τ , or \mathfrak{B}_τ is a $\tau\sim$ -**embrace** of \mathfrak{A}_τ ; symbolically $\mathfrak{A}_\tau \triangleleft \mathfrak{B}_\tau$. If \mathfrak{A}_τ is packed in \mathfrak{B}_τ and $(\mathfrak{A}_\tau, \sim)$, regarded as a partial algebra w.r.t. \sim , is a relative subalgebra of a (full) algebra $(\mathfrak{B}_\tau, \sim)$ (in the sense of G. Grätzer, Universal Algebra, § 13), we also say that $(\mathfrak{A}_\tau, \sim)$ is **packed in** $(\mathfrak{B}_\tau, \sim)$, denoting this by $(\mathfrak{A}_\tau, \sim) \triangleleft (\mathfrak{B}_\tau, \sim)$.

Proposition 4

If $\mathfrak{A}_\tau \triangleleft \mathfrak{B}_\tau$, then \mathfrak{B}_τ is generated by Heyting operations from $|\mathfrak{A}_\tau| \cup \{\sim_\tau\}$.

Two (important) examples

- 1 Given a τ -expansion \mathfrak{A}_τ , \mathfrak{A}_τ^Δ is a $\tau\sim$ -embrace of the former.
- 2 Let Z be a 2-element Heyting algebra. Then $\mathcal{E}_Z = \{(\mathbf{0}, \mathbf{1}), (\mathbf{1}, \mathbf{1})\}$. Accordingly, we have two choices to define \sim :

x	$\sim x$
0	1
1	0

x	$\sim x$
0	1
1	1

Proposition 5

Any nontrivial subvariety of K_τ contains Z_τ .

Proposition 6

Any finite \mathfrak{A}_τ is packed in itself. Hence, Z_τ is packed in itself.

Expansions of Int

Extending \mathcal{L}_a

- The atomic formulas — $Atom = Var \cup \{\tau\}$
- \mathcal{L}_τ -formulas $(A, B, \dots) := A \in Atom | (A \wedge B) | (A \vee B) | (A \rightarrow B) | \neg A$
- $\mathcal{L}_{\tau\sim}$ -formulas $(\alpha, \beta, \dots) :=$
 $\alpha \in Atom | (\alpha \wedge \beta) | (\alpha \vee \beta) | (\alpha \rightarrow \beta) | \neg \alpha | \sim \alpha$

Expanding Int

- \mathbf{Int}_τ denotes intuitionistic propositional logic in \mathcal{L}_τ
- $\mathbf{Int}_{\tau\sim}$ denotes intuitionistic logic in $\mathcal{L}_{\tau\sim}$ plus the following axioms:
 - 1 $\sim p_0 \leftrightarrow (p_0 \rightarrow \tau) \wedge \sim \tau,$
 - 2 $(\sim \tau \rightarrow \tau) \rightarrow \tau,$
 - 3 $\sim \tau \rightarrow (p_0 \vee (p_0 \rightarrow \tau)),$
 - 4 $\tau \rightarrow \sim \tau.$

Proposition 7

$\mathbf{Int}_{\tau\sim} \vdash \alpha \iff \alpha$ is valid in any $\tau\sim$ -expansion.

Proposition 8

For any \mathcal{L}_τ -formulas A and B ,

$$\mathbf{Int}_{\tau\sim} + A \vdash B \iff \mathbf{Int}_\tau + A \vdash B.$$

Corollary 8.1

Any variety of τ -expansions is generated by those algebras of the variety which are τ -reducts of a $\tau\sim$ -embraces.

Remark

In the proof of Corollary 8.1, we use that any variety of τ -expansions contains τ -reducts of some $\tau\sim$ -embraces, for example, Z_τ .

Corollary 8.2

Any τ -expansion \mathfrak{A}_τ is embedded into the τ -reduct \mathfrak{B}_τ of $\tau\sim$ -expansion $(\mathfrak{B}_\tau, \sim)$ so that $\mathbf{HSP}(\mathfrak{A}_\tau) = \mathbf{HSP}(\mathfrak{B}_\tau)$. Hence, for any \mathfrak{A}_τ , there is a \mathfrak{B}_τ such that $\mathfrak{A}_\tau \triangleleft \mathfrak{B}_\tau$ and $\mathbf{HSP}(\mathfrak{A}_\tau) = \mathbf{HSP}(\mathfrak{B}_\tau)$.

Conclusive Step

Propositin 9

Let a Heyting algebra \mathfrak{A} be a subalgebra of \mathfrak{B} and let $a \in \mathfrak{A}$. Assume that a is enriched in \mathfrak{B} by and an element b . If \mathfrak{B} is generated by $|\mathfrak{A}| \cup \{b\}$, then \mathfrak{B} is isomorphic to \mathfrak{A}_a^Δ .

Using Corollary 8.2 and Proposition 9, we obtain the following.

Proposition 10

Let \mathfrak{A} be a Heyting algebra and $a \in \mathfrak{A}$. Then $\text{HSP}(\mathfrak{A}) = \text{HSP}(\mathfrak{A}_a^\Delta)$.

Theorem

For any Heyting algebra \mathfrak{A} , $\text{HSP}(\mathfrak{A}) = \text{HSP}(\vec{\mathfrak{A}})$.

Thank You