

RELATIONAL SEMANTICS FOR DISTRIBUTIVE SUBSTRUCTURAL LOGICS

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CONTENTS

1 OUR LOGICS

- Distributive Substructural Logics

2 RELATIONAL SEMANTICS FOR DFL LOGICS

- Relational Semantics for DFL Logics
- Basic results for DFL-frame
- Contracts with other relational semantics
- Int-frame vs. DFL_{ceW} -frame

3 DESCRIPTION

- General DFL-frame
- Priestley-type Duality
- Topological Characterization of descriptive DFL-frame

4 FUTURE WORK

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- Distributive Substructural Logics

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- 2 RELATIONAL SEMANTICS FOR DFL LOGICS
 - Relational Semantics for DFL Logics
 - Basic results for DFL-frame
 - Contracts with other relational semantics
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DISTRIBUTIVE SUBSTRUCTURAL LOGICS

$$\text{Language} = \begin{cases} p, q, r, \dots & \text{Propositional variables} \\ \mathbf{t}, \mathbf{f}, \top, \perp & \text{Constants} \\ \vee, \wedge, \circ, \backslash, / & \text{Logical connectives} \end{cases}$$

DFL = LJ - Structural rules + Distributivity
Contraction, Exchange, Left-(Right-)weakening

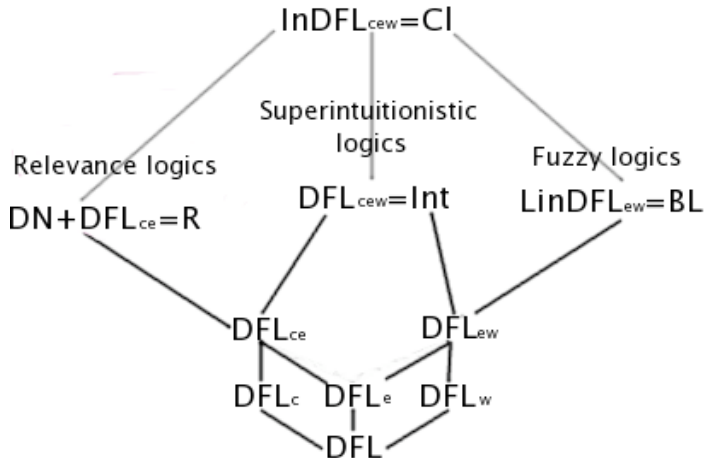
DFL : the set of provable formulas in DFL

L is a DFL logic, if **L** is an extension of **DFL**.

BASIC DFL LOGICS

DFL, **DFL_c**, **DFL_e**, **DFL_w**, **DFL_{ce}**, **DFL_{ew}**, **DFL_{cew}**

IMPORTANT CLASSES OF DFL LOGICS



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RELATIONAL SEMANTICS FOR DFL LOGICS

DEFINITION

A tuple $\mathfrak{F} = \langle W, W_t, W_f, R_o \rangle$ is a DFL-frame, if \mathfrak{F} satisfies the following.

- ① $R_o(w, t_1, w)$ and $R_o(w, w, t_2)$ for some $t_1, t_2 \in W_t$.
- ② If $R_o(w, v, u)$, $w \preceq w'$, $v' \preceq v$ and $u' \preceq u$, then $R_o(w', v', u')$.
- ③ $R_o(w, x, s)$ and $R_o(x, v, u)$ for some $x \in W$, if and only if $R_o(w, v, y)$ and $R_o(y, u, s)$ for some $y \in W$.
- ④ If $w \in W_t$ and $w \preceq w'$, then $w' \in W_t$.
- ⑤ If $w \in W_f$ and $w \preceq w'$, then $w' \in W_f$.

Abbreviation: $w \preceq w' \iff R_o(w', t, w)$ or $R_o(w', w, t)$, for some $t \in W_t$.

$Up(W)$: the set of all subsets of W upward closed under \preceq .

CORRESPONDENCE FOR STRUCTURAL RULES

PROPOSITION

- Contraction $\iff \left\{ \begin{array}{l} \forall w \in W [R_o(w, w, w)] \\ \text{or equivalently} \\ \forall w, v \in W [w \preceq v \Rightarrow R_o(v, w, w)] \end{array} \right.$
- Exchange $\iff \forall w, v, u \in W [R_o(w, v, u) \Rightarrow R_o(w, u, v)]$
- Left weakening $\iff W_t = W$
- Right weakening $\iff W_f = \emptyset$

KRIPKE COMPLETENESS FOR BASIC DFL LOGICS

THEOREM

All basic DFL logics are Kripke complete.

We have already proved several other DFL logics.

CONTRACTS WITH OTHER RELATIONAL SEMANTICS

- Kripke frame for Intuitionistic logic \iff DFL_{cew} -frame
- Routley-Meyer semantics without negation \iff DFL_{ce} -frame without negation

KRIPKE FRAME FOR INTUITIONISTIC LOGIC

DEFINITION (INT-FRAME)

A tuple $\mathfrak{F}_{Int} = \langle W_{Int}, R_{Int} \rangle$ is a Kripke frame for **Int**, if \mathfrak{F}_{Int} satisfies the following.

- 1 For any $w \in W_{Int}$, $wR_{Int}w$.
- 2 If $wR_{Int}v$ and $vR_{Int}u$, then $wR_{Int}u$.

A valuation is a function from the set of propositional variables to $UpR(W_{Int})$.

HOW CAN WE CONSIDER DFL_{cew} -FRAME AS INT-FRAME?

DFL_{cew} -frame: $\mathfrak{F} = \langle W, W_t, W_f, R_o \rangle$

$\Updownarrow?$

Int-frame: $\mathfrak{F}_{Int} = \langle W_{Int}, R_{Int} \rangle$

INT-FRAME VS. DFL_{cew} -FRAME

DFL_{cew} -frame \Rightarrow Int-frame

$$R_o(w, v, u) \Rightarrow v \preceq w$$

- R_o -reflexivity $\Rightarrow R_{Int}$ -reflexivity
- R_o -transitivity $\Rightarrow R_{Int}$ -transitivity

INT-FRAME VS. DFL_{cew} -FRAME

DFL_{cew} -frame \Rightarrow Int-frame

$$R_o(w, v, u) \Rightarrow v \preceq w$$

- R_o -reflexivity $\Rightarrow R_{Int}$ -reflexivity
- R_o -transitivity $\Rightarrow R_{Int}$ -transitivity

INT-FRAME VS. DFL_{cew} -FRAME

Int-frame \Rightarrow DFL_{cew} -frame

$R_o(v, w, u)$, if $wR_{Int}v$ and $uR_{Int}v$.

- R_{Int} -reflexivity \Rightarrow R_o -reflexivity and R_o -idempotency
- R_{Int} -transitivity \Rightarrow R_o -transitivity

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GENERAL DFL-FRAME

DEFINITION

A tuple $\mathfrak{G} = \langle \mathfrak{F}, A \rangle$ is a general DFL-frame, if \mathfrak{F} is a DFL-frame and A satisfies the following.

- ① A is a subset of $Up(W)$.
- ② W_t, W_f, W, \emptyset are included in A .
- ③ A is closed under $\cup, \cap, *, \downarrow$ and \downarrow .

$$X * Y := \{w \in W \mid R_o(w, v, u), v \in X \text{ and } u \in Y, \text{ for some } v, u \in W\}$$

$$X \downarrow Y := \{w \in W \mid R_o(u, v, w), v \in X \Rightarrow u \in Y, \text{ for any } v, u \in W\}$$

$$Y \downarrow X := \{w \in W \mid R_o(u, w, v), v \in X \Rightarrow u \in Y, \text{ for any } v, u \in W\}$$

Anti-symmetry: $w \preceq v$ and $v \preceq w$ only if $w = v$.

PRIESTLEY-TYPE DUALITY

DEFINITION (DUAL ALGEBRA)

Given a general L-frame $\mathfrak{G} = \langle W, W_t, W_f, R_o, A \rangle$, the tuple $\mathfrak{G}^* = \langle A, \cup, \cap, *, \searrow, \swarrow, W_t, W_f, W, \emptyset \rangle$ is *the dual algebra*.

DEFINITION (DUAL FRAME)

Given a dual algebra $\mathfrak{G}^* = \langle A, \cup, \cap, *, \searrow, \swarrow, W_t, W_f, W, \emptyset \rangle$, the tuple $(\mathfrak{G}^*)_* = \langle Pf(A), Pf_{W_t}(A), Pf_{W_f}(A), R_*, \widehat{A} \rangle$ is *the dual frame*

$$R_*(F_1, F_2, F_3) \iff \forall X, Y \in A [X \in F_2 \text{ and } Y \in F_3 \Rightarrow X * Y \in F_1]$$

$$\widehat{A} := \{\widehat{X} \mid X \in A\}$$

$$\widehat{X} := \{F \in Pf(A) \mid X \in F\}$$

PRIESTLEY-TYPE DUALITY

DEFINITION (DUAL ALGEBRA)

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$$\hat{A} := \{ \hat{X} \mid X \in A \}$$

$$\hat{X} := \{ F \in Pf(A) \mid X \in F \}$$

TOPOLOGICAL CHARACTERIZATION OF DESCRIPTIVE DFL-FRAME

DEFINITION

A general L-frame \mathfrak{G} is *descriptive*, if \mathfrak{G} is isomorphic to $(\mathfrak{G}^*)_*$.

THEOREM

A general L-frame \mathfrak{G} is descriptive, if and only if \mathfrak{G} satisfies R_o -tightness and Compactness.

R_o -tightness :

$$R_o(w, v, u) \iff \forall X, Y \in A [v \in X, u \in Y \Rightarrow w \in X * Y]$$

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FUTURE WORK

- Filtration or Finite model property
- Sahlqvist-type Theorems
- etc