# Countable structures with few reducts 

Dugald Macpherson (University of Leeds)

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If $\mathcal{M}^{\prime}$ and $\mathcal{M}$ have the same domain, $\mathcal{M}^{\prime}$ is a group-reduct if $\operatorname{Aut}\left(\mathcal{M}^{\prime}\right) \geq \operatorname{Aut}(\mathcal{M})$, is proper if $\operatorname{Aut}\left(\mathcal{M}^{\prime}\right)>\operatorname{Aut}(\mathcal{M})$, trivial if $\operatorname{Aut}(\mathcal{M})=\operatorname{Sym}(M)$.

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Basic Problem: Find examples of structures with few/no proper non-trivial reducts (in either or both senses).

## Overview of Talk

- Reducts of $\omega$-categorical structures.
- Non- $\omega$-categorical example with no proper non-trivial reducts (both senses).
- Other possible examples (e.g. for strongly minimal sets).

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Remark. There are many (abstractly) maximal subgroups of $\operatorname{Sym}(\mathbb{N})$, e.g. stabilisers of ultrafilters on $\mathbb{N}$ (so $2^{2^{\aleph_{0}}}$ up to conjugacy).

Recall: A permutation group $(G, X)$ is primitive if there is no proper non-trivial $G$-invariant equivalence relation on $X$.

## Remarks.

- If $\mathcal{M}$ is $\omega$-categorical, then the two notions of reduct coincide (Ryll-Nardzewski), and we just call them 'reducts'.

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- If $P$ is a finite subset of $M$, then $(M, P)$ has no proper non-trivial reducts (and intransitive maximal-closed subgroups of $\operatorname{Sym}(M)$ all arise like this).

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So we are interested in maximal-closed subgroups of $\operatorname{Sym}(\mathbb{N})$ which act primitively on $\mathbb{N}$.


## Omega-categoricity

Theorem (Cameron, 1976)
The only proper non-trivial reducts of $(\mathbb{Q},<)$ are

- $(\mathbb{Q}, B)$ (ternary betweenness, ' $x$ is between $y$ and $z$ ')
- $(\mathbb{Q}, K)$ (ternary circular order $K$ induced from $<$ )
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BUT (Junker, Ziegler, 2008) $(\mathbb{Q},<0)$ has 116 reducts!

## Omega-categoricity - homogeneous graphs

## Theorem (Thomas, 1991)

(i) The only proper non-trivial reducts of the random graph $(\Gamma, R)$ are

- $(\Gamma, B)$ ( $B$ ternary, the random graph up to anti-isomorphism)
- $(\mathbb{Q}, K)$ ( $K$ ternary, the homogeneous 'two-graph' induced from $R$, a triple satisfying $K$ iff its entries are distinct and it contains an odd number of graph edges)
- $(\mathbb{Q}, S)$ (the above homogeneous two-graph up to anti-isomorphism)
(ii) For $n \geq 3$, the generic $K_{n}$-free graph has no proper non-trivial reducts.
- Bennett (PhD, Rutgers, 1997): Result for random tournament like that for random graph (3 proper non-trivial reducts).
- (Pach, Pinsker, Pluhár, Pongrácz, Szabó) Similar result for generic poset (3 reducts).
- Thomas (1996): Classification of reducts of random hypergraphs.

Key tool (in treatments of such results by Bodirsky, Pinsker, Tsankov, motivated by constraint satisfaction problems (CSPs)):

## Definition

Let $\mathcal{C}$ be a class of finite relational structures with a language including $<$ (interpreted by a total order). Then $\mathcal{C}$ has the Ramsey Property if for every $A, B \in \mathcal{C}$ and positive integer $k$, there is $D \in \mathcal{C}$ such that for every colouring with $k$ colours of the copies of $A$ in $D$, there is a copy of $B$ in $D$ all of whose substructures isomorphic to $A$ have the same colour. In Ramsey notation,

$$
D \rightarrow(B)_{A}^{k}
$$

For all the above structures $\mathcal{M}$ (random graph, random hypergraph, etc.) there is a Fraïssé-homogeneous expansion $\mathcal{M}^{\prime}=(\mathcal{M},<)$ by a total order such that the Age $\left(\mathcal{M}^{\prime}\right)$ (the class of finite structures which embed in $\mathcal{M}^{\prime}$ ) is a Ramsey class.

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Remark. By Kechris-Pestov-Todorcevic, if $\mathcal{C}$ is as above and $\mathcal{M}$ is the Fraïssé limit, then $\operatorname{Aut}(\mathcal{M})$ is extremely amenable: every continuous action of $\operatorname{Aut}(\mathcal{M})$ on a compact Hausdorff space has a fixed point.

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- Cases where permutation group theory (e.g. Jordan groups see later) is applicable.

Conjecture of Thomas (1991): If $\mathcal{M}$ is a Fraïssé-homogeneous structure over a finite relational language, then $\mathcal{M}$ has just finitely many reducts.

## Reducts of non- $\omega$-categorical structures.

## An example.

Let $(T, R)$ be the unique degree 3 graph-theoretic tree. Two rays (infinite one-way paths) are equivalent if they have infinitely many common vertices. The equivalence classes are called ends. Let $M^{+}$ be the set of ends of $(T, R)$.

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Define $D(x, y ; z, w)$ to hold on $M^{+}$iff one of

- $x=y \wedge x \neq z \wedge x \neq w$
- $z=w \wedge x \neq z \wedge y \neq z$
- $x, y, z, w$ are distinct, and there are rays $\hat{x} \in x, \hat{y} \in y$ etc. such that $\hat{x} \cup \hat{y}$ and $\hat{z} \cup \hat{w}$ are disjoint two-way infinite paths.


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Let $M$ be a countable dense subset of $M^{+}$(i.e. for each $a \in T$ there are $x, y, z \in M$ and rays $\hat{x} \in x, \hat{y} \in y$ and $\hat{z} \in z$ such that $\hat{x} \cup \hat{y}, \hat{y} \cup \hat{z}$, and $\hat{x} \cup \hat{z}$ are all two-way infinite paths through $a$ ). Put $\mathcal{M}:=(M, D)$.


## Observations on ( $M, D$ ).

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- Let $K$ be a countable model of $T h\left(\mathbb{Q}_{2}\right)$ with valuation ring $\mathcal{O}$ having maximal ideal $\mathcal{M}$. On $\mathrm{PG}_{1}(K)$ (viewed as $K \cup\{\infty\}$ ), define $D(x, y ; z, w)$ to hold if and only if the cross-ratio $\frac{(x-z)(y-w)}{(x-w)(y-z)} \in 1+\mathcal{M}$. Then

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- 'D-relations' were axiomatised by Adeleke and Neumann.

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Problem: Find other kinds of non- $\omega$-categorical structures with no proper non-trivial reducts of either kind.

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Guess: The above example ( $M, D$ ) still works, starting from a higher degree tree.

## Automorphisms of $(M, D)$.

If $a \in T$, and $x \in M$ there is a unique ray $x_{a} \in x$ starting at $a$. There is an equivalence relation $E_{a}$ on $M$ : put $E_{a}(x, y)$ iff $x_{a}$ and $y_{a}$ have a common edge of $T$. An $E_{a}$-class is called a cone (at $a$ ). Now:

- for each $a \in T$ there are three $E_{a}$-classes, with $\operatorname{Aut}(\mathcal{M})_{a}$ inducing $S_{3}$ on them;


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- if $U$ is a cone, then $\operatorname{Aut}(U, D)$ is transitive, and any automorphism of $(U, D)$ can be extended by $\mathrm{id}_{M \backslash U}$ to an aut. of $(M, D)$.
- if $A \subset M$ is finite, then $\operatorname{Aut}(M)_{(A)}$ (the pointwise stabiliser of $A$ ) has no finite orbits on $M \backslash A$; hence $\operatorname{acl}(A)=A$.


## Permutation groups.

## Definition

Let $(G, X)$ be a permutation group (group $G$ acting faithfully on $X)$.

- If $A \subset X$ with $|A|>1$, then $A$ is a Jordan set if $G_{(X \backslash A)}$ is transitive on $A$. It is proper if $A \neq X$, and if $|X \backslash A|=n \in \mathbb{N}$, then $(G, X)$ is not $(n+1)$-transitive.
- A Jordan group is a transitive permutation group with a proper Jordan set.

Adeleke, M, Neumann $(1995,1996)$ Structure theorem for primitive Jordan permutation groups. In particular,

## Theorem

Let $G$ be a 3-transitive but not highly transitive Jordan permutation group on an infinite set $X$. Then $G$ preserves on $X$ one of

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Steiner $k$-system on $X$ : collection of 'blocks' (subsets of $X$ all of the same size $>k$ ) such that any $k$ elements of $X$ lie on a unique block.

Sketch proof (that ( $M, D$ ) has no proper non-trivial group-reducts).

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Apply above classification. Rule out all but the $D$-relation case, and show that if $H$ preserves a $D$-relation $D^{\prime}$ on $M$ then it has the same cones as $(M, D)$, so $D^{\prime}=D$.

For example, suppose $H$ preserves a Steiner $k$-system on $M$. Let $a_{1}, \ldots, a_{k} \in M$ be distinct, and let $l$ be the block containing $a_{1}, \ldots, a_{k}$.

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Let $b_{k} \in M$ not lie on $l$, and let $m$ be the block through $a_{1}, \ldots, a_{k-1}, b_{k}$. Choose $a_{k+1} \notin\left\{a_{1}, \ldots, a_{k}\right\}$ on $l$ and $b_{k+1} \notin\left\{a_{1}, \ldots, a_{k-1}, b_{k}\right\}$ on $m$.

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Now $l$ is the unique block containing $a_{1}, \ldots, a_{k-2}, a_{k}, a_{k+1}$, and $m$ the unique block containing $a_{1}, \ldots, a_{k-2}, b_{k}, b_{k+1}$, and these blocks have common points $a_{1}, \ldots, a_{k-1}$.

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Put $A:=\left\{a_{1}, \ldots, a_{k-2}, a_{k}, a_{k+1}, b_{k}, b_{k+1}\right\}$. Then $a_{k-1}$ lies in a singleton orbit of $H_{(A)}$. This is impossible as $a_{k-1} \notin A$ and $H \geq G$.

To show ( $M, D$ ) has no proper non-trivial definable reducts. Let $\mathcal{M}^{\prime}$ be a non-trivial definable reduct of $(M, D)$.

Claim: It suffices to show that the set of cones of $(M, D)$ is uniformly definable in $(M, D)$, i.e. there are formulas $\phi\left(x, y_{1}, \ldots, y_{k}\right)$ and $\psi\left(y_{1}, \ldots, y_{k}\right)$ over $\emptyset$ (in the language of $\mathcal{M}^{\prime}$ ) such that the set of cones of $(M, D)$ is exactly

$$
\left\{\phi\left(M^{\prime}, \bar{a}\right): \bar{a} \in M^{k}, \mathcal{M}^{\prime} \models \psi(\bar{a})\right\} .
$$

Proof: Easy to check that the cones determine $D$.

Step 1. Show that some cone of $(M, D)$ is definable in $\mathcal{M}^{\prime}$, by some formula $\phi(x, \bar{a})$. Hence, all cones are $\mathcal{M}^{\prime}$-definable via $\phi$ (as $\operatorname{Aut}\left(\mathcal{M}^{\prime}\right) \geq \operatorname{Aut}(M, D)$, which is transitive on the set of cones). Now aim to modify $\phi$, and find $\psi$.

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Step 3. Let $\psi_{2}(\bar{y})$ express also that the complement of the set $\phi(M, \bar{y})$ has the form $\phi\left(M, \overline{y^{\prime}}\right)$ for some $\bar{y}^{\prime}$. (Recall that the complement of a cone in $(M, D)$ is also a cone.)

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Step 4. Let $\psi_{3}(\bar{y})$ express a consequence of the Jordan property of cones. Namely,

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\left.\forall \bar{y}^{\prime}\left(\bigwedge_{i=1}^{n} \neg \phi\left(y_{i}^{\prime}, \bar{y}\right) \rightarrow\left(\left(\phi\left(M, \bar{y}^{\prime}\right) \supseteq \phi(M, \bar{y})\right) \vee \phi\left(M, \bar{y}^{\prime}\right) \cap \phi(M, \bar{y})=\emptyset\right)\right)\right)
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Step 5. Reduce to $l(\bar{y})=4$. Reduce to case when $\phi\left(M, \bar{a}^{\prime}\right)$ is a cone or union of two cones at adjacent nodes. Finish.

## More on reducts

Partially order definable reducts of $\mathcal{M}$, putting $\mathcal{M}_{1} \leq \mathcal{M}_{2}$ iff $\mathcal{M}_{2}$ is $\emptyset$-definable in $\mathcal{M}_{1}$ (factoring out equi-definability). Partially order group-reducts by group inclusion.

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Lemma
Let $\mathcal{M}$ be saturated.
(i) The partial order of definable reducts of $\mathcal{M}$ embeds into the partial order of group-reducts of $\mathcal{M}$.
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Proof. Easy compactness and saturation argument.

## Reducts of strongly minimal sets

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First try. Let $(T, R)$ be the degree 3 tree (strongly minimal, disintegrated, i.e. $\operatorname{acl}(A)=\bigcup(\operatorname{acl}(a): a \in A)$ for any $A)$. The 'distance 2' graph $T^{(2)}$ is a proper non-trivial definable and group reduct. $T^{(2)}$ is the disjoint union of two graphs, each built from copies of $K_{3}$ in a treelike way, three copies of $K_{3}$ containing each vertex. There is also a (non-definable) group-reduct, the equivalence relation with 2 classes corresponding to 'even distance apart'. Any other reducts?

Second try. Let $k, l \in \mathbb{N}$ with $k \geq 2$ and $l \geq 3$. Let $\Gamma_{k, l}$ be the graph consisting of copies of $K_{k+1}$ put together in a treelike way, with $l$ copies of $K_{k+1}$ containing each vertex.

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Theorem (Bodirsky, M)
$\mathcal{M}:=\Gamma_{k, l}$ has no proper non-trivial definable reducts.

## Sketch Proof.

1. Any definable reduct $\mathcal{M}^{\prime}$ is strongly minimal and disintegrated, so if $\mathcal{M}^{\prime}$ is non-trivial then $\left|\operatorname{acl}_{\mathcal{M}^{\prime}}(a)\right|>1$ for any $a$. Hence in $\mathcal{M}^{\prime}$, some set $\phi(M, a)$ has finite size greater than 1 . By distance transitivity, can assume there are $1 \leq n_{1}<\ldots<n_{t} \in \mathbb{N}$ such that $\phi(x, y)$ is ' $d(x, y) \in\left\{n_{1}, \ldots, n_{t}\right\}$ '.

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2. Put $n:=n_{t}$. Show that for any vertices $x, y$,

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d(x, y) \leq 2 n \Leftrightarrow \mathcal{M}^{\prime} \mid=\exists z(\phi(x, z) \wedge \phi(y, z)) .
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Thus the balls $B_{2 n}(x)$ are uniformly definable.
3. Show there is some $\gamma \in \mathbb{N}$ such that

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x, y \text { are adjacent in } \mathcal{M} \Leftrightarrow B_{2 n}(x) \cap B_{2 n}(y)<\gamma
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Non-disintegrated locally modular strongly minimal sets. For finite fields $F$, vector spaces over $F$ (and projective and affine spaces over $F$ ) are $\omega$-categorical, and reducts can be handled by the Cherlin-Harrington-Lachlan and Zilber work.

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If $F$ has characteristic $p$, there are proper reducts by viewing $V$ as over the prime subfield.

## Further questions

1. Define $f$ on $\mathbb{Q}$ by $f(x, y, z):=x-y+z$. Does $(\mathbb{Q}, f)$ have any proper non-trivial definable reducts?

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3. Does $\operatorname{Sym}(\mathbb{N})$ have any countable maximal-closed subgroups?
4. Is every closed proper subgroup of $\operatorname{Sym}(\mathbb{N})$ contained in a maximal-closed subgroup of $\operatorname{Sym}(\mathbb{N})$ ?

