## On o-minimal MV-chains

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Algebraic Semantics for Uncertainty and Vagueness

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- 2 MV-chains and O-minimality
- 3 Perfect MV-chains
- 4 Local MV-chains of finite rank
- 5 Imaginary Elements

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- A first-order theory T is said to be *o-minimal* if every model of T is o-minimal.

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- (iii) a real closed field in the language  $L = \langle +, \cdot, 0, 1, < \rangle$ .

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- Indeed, the only definable non-trivial subgroup of an o-minimal ordered abelian group  ${\bf G}$  is  ${\bf G}$  itself.
- Moreover, in any o-minimal ordered abelian group it is possible to define a divisible group.
- We have that an ordered abelian group is divisible IFF it is o-minimal IFF its theory admits QE [Pillay, Steinhorn (1986); Lenski (1989)].

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- The theory of the ordered group of integers in the language  $\langle +, -, 0, <, \{P_n\}\rangle$ , with  $n = 2, 3, \ldots$  has QE but is not o-minimal. Indeed, the formula  $\exists y \ 2y = x$  defines an infinite union of intervals.

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- The theory of BL-chains representable as an infinite ordinal sum ⊕<sub>i∈I</sub> A<sub>i</sub> of divisible MV-chains, with a densely ordered index set *I* with a minimum and without a maximum has QE in the language of BL-chains. However, the set of idempotents is definable [Cortonesi, M., Montagna (2010)].

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- Every divisible MV-chain is o-minimal...
- but not every o-minimal MV-chain is divisible.

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By an unnested atomic formula in L we mean one of the following formulas:
(i) x = y, (x < y);</li>
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- A formula is called unnested if all its atomic subformulas are unnested.
- For a first-order language  $L = \langle <, f_1, \dots, f_n, c_1, \dots, c_m \rangle$ , every formula is equivalent to an unnested formula.

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- T<sub>1</sub> is *interpretable* into T<sub>2</sub> (with parameters) if there exists an L<sub>2</sub>-formula χ(z), and for every M<sub>1</sub> ⊨ T<sub>1</sub> there exists a M<sub>2</sub> ⊨ T<sub>2</sub> (unique up to isomorphism) such that:
  - (i) there exists a bijection  $h_{M_1}: M_1 \to \{a \mid M_2 \models \chi(a)\}$  from the domain of  $M_1$  into the set defined by the L<sub>2</sub>-formula  $\chi(z)$  over the domain of  $M_2$ ;

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  - (i) there exists a bijection  $h_{M_1}: M_1 \to \{a \mid M_2 \models \chi(a)\}$  from the domain of  $M_1$  into the set defined by the L<sub>2</sub>-formula  $\chi(z)$  over the domain of  $M_2$ ;
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## Theorem

Let  $T_1$  and  $T_2$  be two theories in the languages  $L_1$  and  $L_2$ , respectively. Suppose that  $T_1$  is interpretable in  $T_2$ , and  $T_2$  is o-minimal. Then,  $T_1$  is o-minimal as well.

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# Corollary

The theory of divisible MV-chain is interpretable in the theory of ordered divisible abelian groups, therefore it is o-minimal.

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The theory of divisible MV-chain is interpretable in the theory of ordered divisible abelian groups, therefore it is o-minimal.

Are all o-minimal MV-chains divisible?

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- If such an *n* exists, then ord(x) = n, while if it does not exists,  $ord(x) = \infty$ .
- An MV algebra is called *perfect* if for every x ≠ 0, ord(x) = ∞ if and only if ord(¬(x)) < ∞.</li>

• Let  $\mathbb{Z} \times G$ , where G is an ordered abelian group, be an ordered abelian group equipped with the lexicographic order.

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- Every perfect MV chain  $\mathbf{A} \cong \Gamma(\mathbb{Z} \times \mathbf{G}, (1, 0))$  [Di Nola, Lettieri (1994)].

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### Theorem

Every Perfect MV-chain is semidivisible IFF it is o-minimal.

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- A local MV-algebra A of rank n is radical retractive if A/Rad(A) is a subalgebra of A.

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## Theorem

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- An equivalence formula of A is a formula φ(x̄, ȳ) of L, without parameters, such that the relation {(a, b) : A ⊨ φ(ā, b̄)} is a non-empty equivalence relation E<sub>φ</sub>.

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- L-structure **A** has elimination of imaginaries if for every equivalence formula  $\theta(\overline{x}, \overline{y})$  of **A** and each tuple  $\overline{a}$  in **A** there is a formula  $\phi(\overline{x}, \overline{y})$  of L such that the equivalence class  $\overline{a}/\theta$  of  $\overline{a}$  can be written as  $\phi(A^n, \overline{b})$  for some unique tuple  $\overline{b}$  from **A**.

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- We say that **A** has *uniform elimination of imaginaries* if the same holds, except that  $\phi$  depends only on  $\theta$  and not on  $\overline{a}$ .

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• Every o-minimal structure with definable Skolem functions and at least two constant elements has uniform elimination of imaginaries [Hodges (1993)].

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- Every o-minimal structure with definable Skolem functions and at least two constant elements has uniform elimination of imaginaries [Hodges (1993)].
- Recall that a theory T has definable Skolem functions if for every formula  $\phi(\overline{x}, y)$ , with  $\overline{x}$  not empty, there is a formula  $\psi(\overline{x}, y)$  such that

$$\mathsf{T} \vdash \forall \overline{\mathsf{x}} \ (\exists \mathsf{y} \ \phi(\overline{\mathsf{x}},\mathsf{y}) \to (\exists_{=1}\mathsf{y} \ \psi(\overline{\mathsf{x}},\mathsf{y}) \land \forall \mathsf{y} \ (\psi(\overline{\mathsf{x}},\mathsf{y}) \to \phi(\overline{\mathsf{x}},\mathsf{y})))).$$

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  - every MV-chain A can be embedded into a divisible one B such that there is no automorphism of B fixing A other than the identity.

# **THANKS!**