## On o-Minimal MV-CHAINS

## Enrico Marchioni

Artificial Intelligence Research Institute (IIIA - CSIC), Spain enrico@iiia.csic.es

Algebraic Semantics for Uncertainty and Vagueness
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(1) O-Minimality
(2) MV-chains and O-minimality
(3) Perfect MV-chains
(4) Local MV-chains of finite rank
(5) Imaginary Elements

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- A first-order theory $T$ is said to be o-minimal if every model of $T$ is o-minimal.


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(iii) a real closed field in the language $L=\langle+, \cdot, 0,1,<\rangle$.

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- Any o-minimal ordered abelian group is also divisible [Pillay, Steinhorn (1986)].
- Indeed, the only definable non-trivial subgroup of an o-minimal ordered abelian group $\mathbf{G}$ is $\mathbf{G}$ itself.
- Moreover, in any o-minimal ordered abelian group it is possible to define a divisible group.
- We have that an ordered abelian group is divisible IFF it is o-minimal IFF its theory admits QE [Pillay, Steinhorn (1986); Lenski (1989)].


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- The theory of BL-chains representable as an infinite ordinal sum $\bigoplus_{i \in I} \mathbf{A}_{i}$ of divisible MV-chains, with a densely ordered index set I with a minimum and without a maximum has QE in the language of BL-chains. However, the set of idempotents is definable [Cortonesi, M., Montagna (2010)].


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- What happens with MV-algebras?
- Can we exploit the characterization for ordered Abelian groups?
- Every divisible MV-chain is o-minimal...
- but not every o-minimal MV-chain is divisible.


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- By an unnested atomic formula in L we mean one of the following formulas:
(i) $x=y, \quad(x<y)$;
(ii) $\quad x=c, \quad(x<c), \quad$ for some constant symbol $c \in L$;
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- A formula is called unnested if all its atomic subformulas are unnested.
- For a first-order language $\mathrm{L}=\left\langle<, f_{1}, \ldots, f_{n}, c_{1}, \ldots, c_{m}\right\rangle$, every formula is equivalent to an unnested formula.


## Interpretation (II)

- Let $\mathrm{T}_{1}$ and $\mathrm{T}_{2}$ be two theories in the the languages $\mathrm{L}_{1}=\left\langle<, f_{1}, \ldots, f_{n}, c_{1}, \ldots, c_{m}\right\rangle$ and $\mathrm{L}_{2}=\left\langle<, f_{1}^{\prime}, \ldots, f_{n^{\prime}}^{\prime}, c_{1}^{\prime}, \ldots, c_{m^{\prime}}^{\prime}\right\rangle$, respectively.


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- $\mathrm{T}_{1}$ is interpretable into $\mathrm{T}_{2}$ (with parameters) if there exists an $\mathrm{L}_{2}$-formula $\chi(z)$, and for every $\mathbf{M}_{1} \models \mathrm{~T}_{1}$ there exists a $\mathbf{M}_{2} \models \mathrm{~T}_{2}$ (unique up to isomorphism) such that:


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\mathbf{M}_{1} \models \varphi(\bar{b}) \text { if and only if } \mathbf{M}_{2} \models \varphi^{\sharp}\left(h_{M_{1}}(\bar{b})\right) .
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- Let $\mathrm{T}_{1}$ and $\mathrm{T}_{2}$ be two theories in the languages $\mathrm{L}_{1}$ and $\mathrm{L}_{2}$, respectively. Suppose that $\mathrm{T}_{1}$ is interpretable in $\mathrm{T}_{2}$. Then for every $\mathbf{M}_{1} \models \mathrm{~T}_{1}$ and for each $\mathrm{L}_{1}$-formula $\varphi(\bar{x})$, there exists an $\mathrm{L}_{2}$-formula $\varphi^{\sharp}(\bar{x})$ so that, for all $\bar{b} \in M_{1}$


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## O-minimality and Interpretation

## Theorem

Let $\mathrm{T}_{1}$ and $\mathrm{T}_{2}$ be two theories in the languages $\mathrm{L}_{1}$ and $\mathrm{L}_{2}$, respectively. Suppose that $\mathrm{T}_{1}$ is interpretable in $\mathrm{T}_{2}$, and $\mathrm{T}_{2}$ is o-minimal. Then, $\mathrm{T}_{1}$ is o-minimal as well.

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## Corollary

The theory of divisible MV-chain is interpretable in the theory of ordered divisible abelian groups, therefore it is o-minimal.

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Are all o-minimal MV-chains divisible?

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- If such an $n$ exists, then $\operatorname{ord}(x)=n$, while if it does not exists, $\operatorname{ord}(x)=\infty$.
- An MV algebra is called perfect if for every $x \neq 0, \operatorname{ord}(x)=\infty$ if and only if $\operatorname{ord}(\neg(x))<\infty$.


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(a, b) \oplus(c, d)= \begin{cases}(a+c, b+d) & a+c<1 \text { or } \\ (1,0) & a+c=1 \text { and } b+d<0 \\ \text { otherwise }\end{cases}
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- $\Gamma(\mathbb{Z} \times \mathbf{G},(1,0))=\langle A, \oplus, \neg,(0,0),(1,0)\rangle$ is a perfect MV-chain
- Every perfect MV chain $\mathbf{A} \cong \Gamma(\mathbb{Z} \times \mathbf{G},(1,0))$ [Di Nola, Lettieri (1994)].


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- An MV-algebra is perfect IFF it satisfies the sentence [Belluce, Di Nola, Gerla (2007)]

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- Semidivisible perfect MV-chains are exactly those chains A such that $\mathbf{A} \cong \Gamma(\mathbb{Z} \times \mathbf{G},(1,0))$, where $\mathbf{G}$ is a divisible ordered abelian group.


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## Theorem

Every Perfect MV-chain is semidivisible IFF it is o-minimal.

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- An MV algebra $\mathbf{A}$ has rank $n$ iff $\mathbf{A} / \operatorname{Rad}(\mathbf{A}) \cong \mathbf{S}_{n}$
- A local MV-algebra $\mathbf{A}$ of rank $n$ is radical retractive if $\mathbf{A} / \operatorname{Rad}(\mathbf{A})$ is a subalgebra of A.


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## Local MV-chains of finite rank (II)

- Let $\mathbb{Z} \times \mathbf{G}$, where $\mathbf{G}$ is an ordered abelian group, be an ordered abelian group equipped with the lexicographic order.
- Let $A=\{x: x \in[(0,0),(n, 0)]\}$, and define over $A$

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\begin{aligned}
&(a, b) \oplus(c, d)= \begin{cases}(a+c, b+d) & a+c<n \text { or } \\
(n, 0) & \begin{array}{l}
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\end{array} \\
\neg(a, b) & =(n-a, 0-b)\end{cases} \\
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- $\Gamma(\mathbb{Z} \times \mathbf{G},(n, 0))=\langle A, \oplus, \neg,(0,0),(n, 0)\rangle$ is a radical retractive local MV-chain of rank $n$.


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- $\Gamma(\mathbb{Z} \times \mathbf{G},(n, 0))=\langle A, \oplus, \neg,(0,0),(n, 0)\rangle$ is a radical retractive local MV-chain of rank $n$.
- Every radical retractive local MV-chain of rank $n \mathbf{A} \cong \Gamma(\mathbb{Z} \times \mathbf{G},(n, 0))$ [Di Nola, Esposito, Gerla (2007)].


## Local MV-chains of finite rank (III)

- Every radical retractive local MV-chain of rank $n$ satisfies [Di Nola, Esposito, Gerla (2007)]

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\forall x\left((2 x=1) \sqcup\left(x^{2}=0\right) \sqcup((n+1) x=1) \sqcap\left(x^{n+1}=0\right)\right) .
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for all $m$, hold in $\mathbf{A}$.

- Semidivisible radical retractive local MV-chain of rank $n$ are exactly those chains A such that $\mathbf{A} \cong \Gamma(\mathbb{Z} \times \mathbf{G},(n, 0))$, where $\mathbf{G}$ is a divisible ordered abelian group.


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## Theorem

Every semidivisible radical retractive local MV-chain of rank $n$ is o-minimal.

## Imaginary Elements (I)

- Let $L$ be a first-order language and $\mathbf{A}$ an $L$-structure.


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- Items of the form $\bar{a} / \phi$, where $\phi$ is an equivalence formula and $\bar{a}$ a tuple, are known as imaginary elements of A.
- L-structure $\mathbf{A}$ has elimination of imaginaries if for every equivalence formula $\theta(\bar{x}, \bar{y})$ of $\mathbf{A}$ and each tuple $\bar{a}$ in $\mathbf{A}$ there is a formula $\phi(\bar{x}, \bar{y})$ of $L$ such that the equivalence class $\bar{a} / \theta$ of $\bar{a}$ can be written as $\phi\left(A^{n}, \bar{b}\right)$ for some unique tuple $\bar{b}$ from $\mathbf{A}$.


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- An equivalence formula of $\mathbf{A}$ is a formula $\phi(\bar{x}, \bar{y})$ of $L$, without parameters, such that the relation $\{(a, b): \mathbf{A} \models \phi(\bar{a}, \bar{b})\}$ is a non-empty equivalence relation $E_{\phi}$.
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- We say that $\mathbf{A}$ has uniform elimination of imaginaries if the same holds, except that $\phi$ depends only on $\theta$ and not on $\bar{a}$.


## Imaginary Elements (II)

Theorem
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- Every o-minimal structure with definable Skolem functions and at least two constant elements has uniform elimination of imaginaries [Hodges (1993)].


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The theory of divisible MV-chains has uniform elimination of imaginaries.

- Every o-minimal structure with definable Skolem functions and at least two constant elements has uniform elimination of imaginaries [Hodges (1993)].
- Recall that a theory T has definable Skolem functions if for every formula $\phi(\bar{x}, y)$, with $\bar{x}$ not empty, there is a formula $\psi(\bar{x}, y)$ such that

$$
\mathrm{T} \vdash \forall \bar{x}(\exists y \phi(\bar{x}, y) \rightarrow(\exists=1 y \psi(\bar{x}, y) \wedge \forall y(\psi(\bar{x}, y) \rightarrow \phi(\bar{x}, y)))) .
$$

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(1) the theory of divisible MV-chains has quantifier elimination;
(2) every MV-chain $\mathbf{A}$ can be embedded into a divisible one $\mathbf{B}$ such that for every $b \in B$ there is a formula $\phi(x)$ (with parameters from $A$ ) such that $\mathbf{B} \models \phi(b)$ and B $\models(\exists \leq n x) \phi(x)$ for some $n$;


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(3) every MV-chain A can be embedded into a divisible one $\mathbf{B}$ such that there is no automorphism of $\mathbf{B}$ fixing $\mathbf{A}$ other than the identity.


## The End

## THANKS!

